

Topology – week 10

Math3061

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© Semester 2, 2022

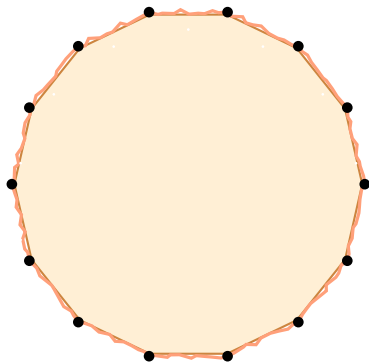
Words for surfaces

A polygonal decomposition for a surface that has **one face** can be encoded in a **word**

Words for surfaces

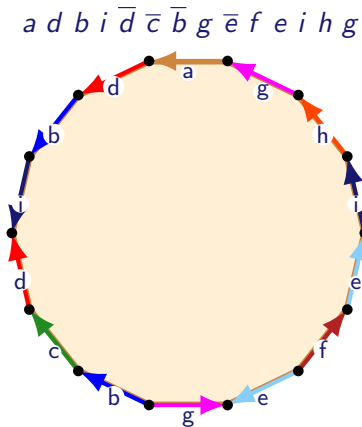
A polygonal decomposition for a surface that has **one face** can be encoded in a **word**

$a d b i \bar{d} \bar{c} \bar{b} g \bar{e} f e i h g$



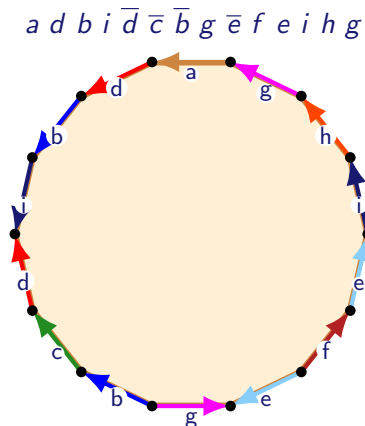
Words for surfaces

A polygonal decomposition for a surface that has **one face** can be encoded in a **word**



Words for surfaces

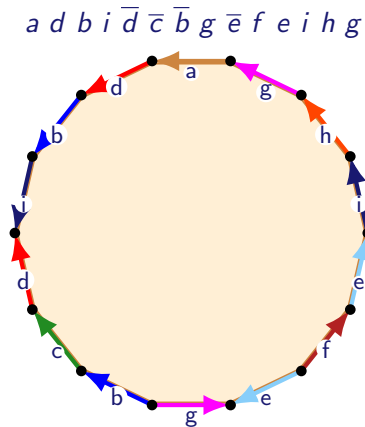
A polygonal decomposition for a surface that has **one face** can be encoded in a **word**



- ▶ write x for an edge pointing **anticlockwise**
- ▶ write \bar{x} for an edge pointing **clockwise**

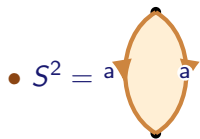
Words for surfaces

A polygonal decomposition for a surface that has **one face** can be encoded in a **word**

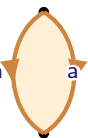


- ▶ write x for an edge pointing **anticlockwise**
- ▶ write \bar{x} for an edge pointing **clockwise**
- ▶ We always read the word in **anticlockwise** order

Words for basic surfaces

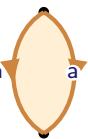


Words for basic surfaces

- $S^2 = a \bar{a}$

 $= a \bar{a}$

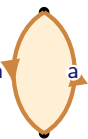
Words for basic surfaces

• $S^2 = a \bar{a}$
 $= a \bar{a}$



The diagram shows a sphere-like shape with two black dots at the top and bottom poles. Two orange curved arrows represent generators: one on the left pointing downwards labeled 'a', and one on the right pointing upwards labeled 'a'.

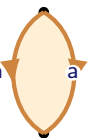
• $\mathbb{P}^2 = a \bar{a}$



The diagram shows a sphere-like shape with two black dots at the top and bottom poles. Two orange curved arrows represent generators: one on the left pointing downwards labeled 'a', and one on the right pointing upwards labeled 'a'.

Words for basic surfaces

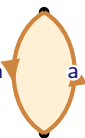
• $S^2 = a \bar{a}$



$= a \bar{a}$

The diagram shows a sphere S^2 represented as a yellow oval with two black dots at the top and bottom poles. Two orange curved arrows, labeled a , represent generators. One arrow starts at the top pole and points left, while the other starts at the bottom pole and points right. Below the diagram, the word $a \bar{a}$ is written in blue.

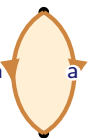
• $\mathbb{P}^2 = a a$



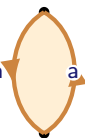
$= a a$

The diagram shows a projective plane \mathbb{P}^2 represented as a yellow oval with two black dots at the top and bottom poles. Two orange curved arrows, both labeled a , represent generators. One arrow starts at the top pole and points left, while the other starts at the bottom pole and points right. Below the diagram, the word $a a$ is written in blue.

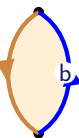
Words for basic surfaces

• $S^2 = a \bar{a}$

 $= a \bar{a}$

The diagram shows a sphere-like shape with two black dots at the top and bottom. Two orange curved arrows represent generators: one on the left pointing downwards and one on the right pointing upwards, both labeled 'a'.

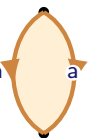
• $\mathbb{P}^2 = a a$

 $= a a$

The diagram shows a sphere-like shape with two black dots at the top and bottom. Two orange curved arrows represent generators: one on the left pointing downwards and one on the right pointing upwards, both labeled 'a'.

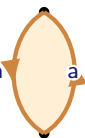
• $\mathbb{D}^2 = a b$


The diagram shows a sphere-like shape with two black dots at the top and bottom. The left orange curved arrow is labeled 'a' and points downwards. The right curved arrow is blue and labeled 'b', pointing upwards.

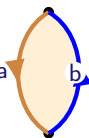
Words for basic surfaces

• $S^2 = a \bar{a}$

 $= a \bar{a}$

The diagram shows a sphere-like shape with two black dots at the top and bottom. Two orange curved arrows represent generators: one on the left pointing downwards and one on the right pointing upwards, both labeled 'a'.

• $\mathbb{P}^2 = a a$

 $= a a$

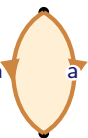
The diagram shows a sphere-like shape with two black dots at the top and bottom. Two orange curved arrows represent generators: one on the left pointing downwards and one on the right pointing upwards, both labeled 'a'.

• $\mathbb{D}^2 = a b$

 $= a b$

The diagram shows a sphere-like shape with two black dots at the top and bottom. Two curved arrows represent generators: one orange arrow on the left pointing downwards labeled 'a', and one blue arrow on the right pointing upwards labeled 'b'.


Words for basic surfaces

• $S^2 = a \bar{a}$
 $= a \bar{a}$



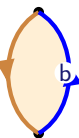
The diagram shows a lens-shaped region with two boundary components, each labeled 'a'. The top boundary is oriented counter-clockwise, and the bottom boundary is oriented clockwise.

• $\mathbb{P}^2 = a a$
 $= a a$



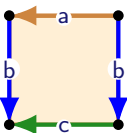
The diagram shows a lens-shaped region with two boundary components, each labeled 'a'. Both the top and bottom boundaries are oriented counter-clockwise.

• $\mathbb{D}^2 = a b$
 $= a b$



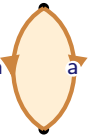
The diagram shows a lens-shaped region with two boundary components. The left boundary is labeled 'a' and oriented counter-clockwise. The right boundary is labeled 'b' and oriented clockwise.


• $\mathbb{A} = b c b$

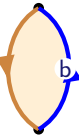


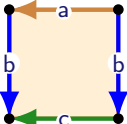
The diagram shows a square region with four boundary components. The top edge is labeled 'a' and oriented counter-clockwise. The bottom edge is labeled 'c' and oriented counter-clockwise. The left and right edges are both labeled 'b' and oriented clockwise.

Words for basic surfaces

• $S^2 = a \overleftarrow{a}$

 $= a \bar{a}$

• $\mathbb{P}^2 = a \overleftarrow{a}$

 $= a a$

• $\mathbb{D}^2 = a \overleftarrow{a}$

 $= a b$

• $\mathbb{A} = b \overleftarrow{b}$

 $= a b \bar{c} \bar{b}$

Words for basic surfaces

• $S^2 = a \overleftarrow{a}$
 $= a \bar{a}$

• $\mathbb{P}^2 = a \overleftarrow{a}$
 $= a a$

• $\mathbb{D}^2 = a \overleftarrow{a} b$
 $= a b$

• $\mathbb{A} = a \overleftarrow{a} b \overleftarrow{b} c \overleftarrow{c}$
 $= a b \bar{c} \bar{b}$

• $\mathbb{M} = a \overleftarrow{a} b \overleftarrow{b} c \overleftarrow{c}$

Words for basic surfaces

• $S^2 = a \overline{a}$
 $= a \bar{a}$

• $\mathbb{P}^2 = a a$


• $\mathbb{D}^2 = a b$

• $\mathbb{A} = a b \bar{c} \bar{b}$
 $= a b \bar{c} \bar{b}$


• $\mathbb{M} = a b \bar{c} b$
 $= a b \bar{c} b$

Words for basic surfaces

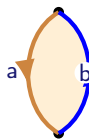
• $S^2 = a \overline{a}$
 $= a \bar{a}$



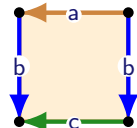
• $\mathbb{P}^2 = a a$



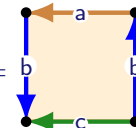
• $\mathbb{D}^2 = a b$



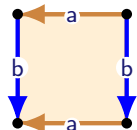
• $\mathbb{A} = a b \bar{c} \bar{b}$
 $= a b \bar{c} \bar{b}$



• $\mathbb{M} = a b \bar{c} b$
 $= a b \bar{c} b$




• $\mathbb{T} = a b a$




Words for basic surfaces

• $S^2 = a \overline{a}$



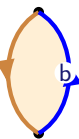
$= a \bar{a}$

• $\mathbb{P}^2 = a a$



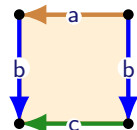
$= a a$

• $\mathbb{D}^2 = a b$



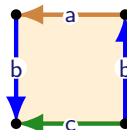
$= a b$

• $\mathbb{A} = a b \bar{c} \bar{b}$



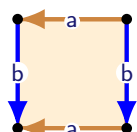
$= a b \bar{c} \bar{b}$

• $\mathbb{M} = a b \bar{c} b$



$= a b \bar{c} b$


• $\mathbb{T} = a b \bar{a} \bar{b}$



$= a b \bar{a} \bar{b}$


Words for basic surfaces

• $S^2 = a \overline{a}$



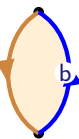
$= a \bar{a}$

• $\mathbb{P}^2 = a a$



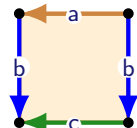
$= a a$

• $\mathbb{D}^2 = a b$



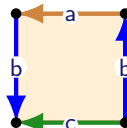
$= a b$

• $\mathbb{A} = a b \bar{c} \bar{b}$



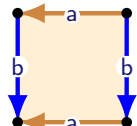
$= a b \bar{c} \bar{b}$

• $\mathbb{M} = a b \bar{c} b$



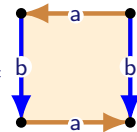
$= a b \bar{c} b$

• $\mathbb{T} = a b \bar{a} \bar{b}$



$= a b \bar{a} \bar{b}$


• $\mathbb{K} = a b \bar{a} \bar{b}$



$= a b \bar{a} \bar{b}$

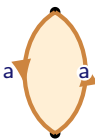
Words for basic surfaces

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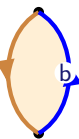
$= a \bar{a}$

• $\mathbb{P}^2 = a a$



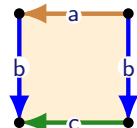
$= a a$

• $\mathbb{D}^2 = a b$



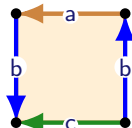
$= a b$

• $\mathbb{A} = a b \bar{c} \bar{b}$



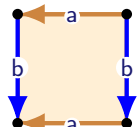
$= a b \bar{c} \bar{b}$

• $\mathbb{M} = a b \bar{c} b$



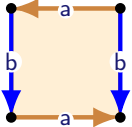
$= a b \bar{c} b$

• $\mathbb{T} = a b \bar{a} \bar{b}$



$= a b \bar{a} \bar{b}$

• $\mathbb{K} = a b a \bar{b}$



$= a b a \bar{b}$

Properties of words

- Words **encode** orientability

- ▶ Orientable: $\dots a \dots \bar{a} \dots$ or $\dots \bar{a} \dots a \dots$

- ▶ Non-orientable: $\dots a \dots a \dots$ or $\dots \bar{a} \dots \bar{a} \dots$

Properties of words

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- Words give a compact and easily readable way of describing surfaces

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Example The following words are all words for the torus \mathbb{T} :

$$\begin{array}{cccc} a b \bar{a} \bar{b} & b \bar{a} \bar{b} a & \bar{a} \bar{b} a b & \bar{b} a b \bar{a} \\ a \bar{b} \bar{a} b & \bar{b} \bar{a} b a & \bar{a} b a \bar{b} & b a \bar{b} \bar{a} \end{array}$$

Properties of words

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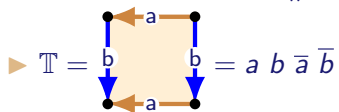
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- The word of a surface can be used to give generators and relations for the first **homotopy group** of the surface — this generalises **independent cycles** and are beyond the scope of this unit

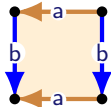
Standard words for closed orientable surfaces

- Connected sums of tori: $\#^t \mathbb{T}$

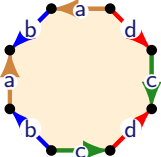


Standard words for closed orientable surfaces

- Connected sums of tori: $\#^t \mathbb{T}$

▶ $\mathbb{T} =$  $= a b \bar{a} \bar{b}$

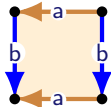
The diagram shows a square with four vertices marked by black dots. The top edge is a horizontal line segment labeled 'a' with an arrow pointing to the left. The right edge is a vertical line segment labeled 'b' with an arrow pointing downwards. The bottom edge is a horizontal line segment labeled 'a' with an arrow pointing to the left. The left edge is a vertical line segment labeled 'b' with an arrow pointing downwards.

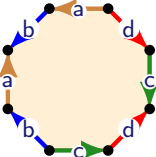
▶ $\#^2 \mathbb{T} =$  $= a b \bar{a} \bar{b} c d \bar{c} \bar{d}$

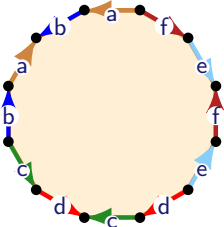
The diagram shows an octagon with eight vertices marked by black dots. The boundary is divided into four pairs of opposite edges, each pair representing a handle. The top edge is labeled 'a' (orange arrow pointing right), the right edge is labeled 'd' (red arrow pointing down), the bottom edge is labeled 'c' (green arrow pointing left), and the left edge is labeled 'b' (blue arrow pointing up). The corresponding reverse edges are labeled 'a-bar' (orange arrow pointing left), 'd-bar' (red arrow pointing up), 'c-bar' (green arrow pointing right), and 'b-bar' (blue arrow pointing down).

Standard words for closed orientable surfaces

- Connected sums of tori: $\#^t \mathbb{T}$

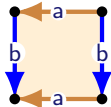
▶ $\mathbb{T} =$  $= a b \bar{a} \bar{b}$

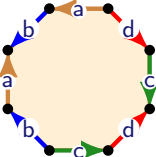
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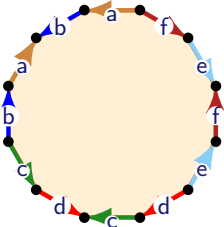
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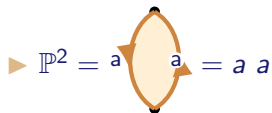
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Words for closed non-orientable surfaces

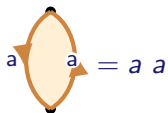
- Connected sums of projective planes $\#^P \mathbb{P}^2$



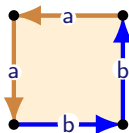
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▶ $\mathbb{P}^2 = a a = a a$




▶ $\#^2 \mathbb{P}^2 = a a b b = a a b b$



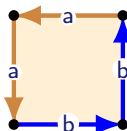
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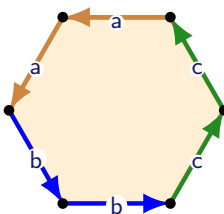
▶ $\mathbb{P}^2 = a \overleftarrow{a} = a a$



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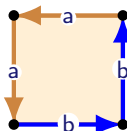
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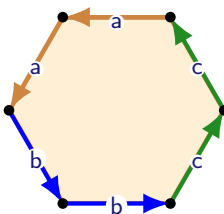
▶ $\mathbb{P}^2 = a a = a a$

A diagram of a projective plane represented as a disk with a boundary loop. The loop is a single line with arrows pointing in opposite directions, labeled 'a' at both ends.

▶ $\#^2 \mathbb{P}^2 = a a b b = a a b b$

A diagram of a connected sum of two projective planes represented as a square. The top and bottom edges are labeled 'a' with arrows pointing outwards. The left and right edges are labeled 'b' with arrows pointing inwards.

▶ $\#^3 \mathbb{P}^2 = a a b b c c = a a b b c c$

A diagram of a connected sum of three projective planes represented as a hexagon. The top and bottom edges are labeled 'a' with arrows pointing outwards. The left and right edges are labeled 'b' with arrows pointing inwards. The two slanted edges are labeled 'c' with arrows pointing outwards.

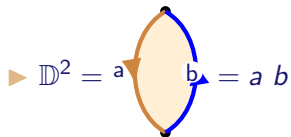
▶ ... $\#^p \mathbb{P}^2 = a_1 a_1 a_2 a_2 \dots a_p a_p$

Standard words for surfaces with boundary

- $\#^d \mathbb{D}^2$

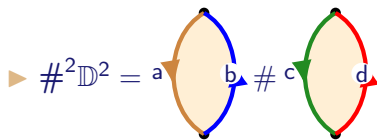
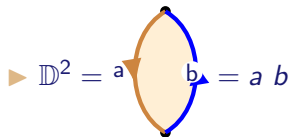
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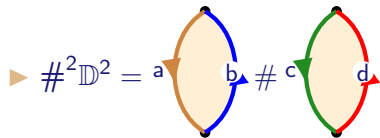
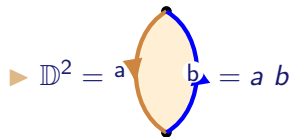
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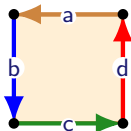


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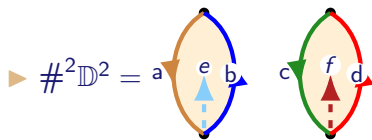
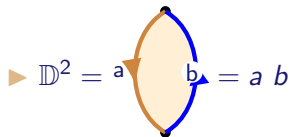


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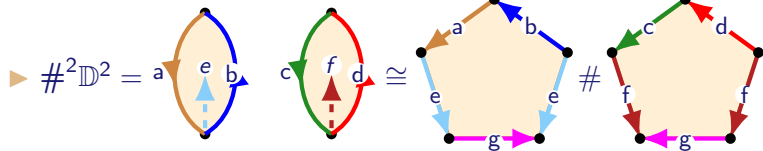
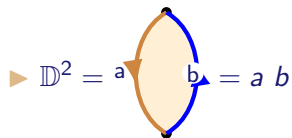
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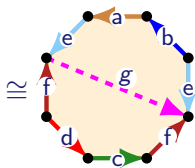
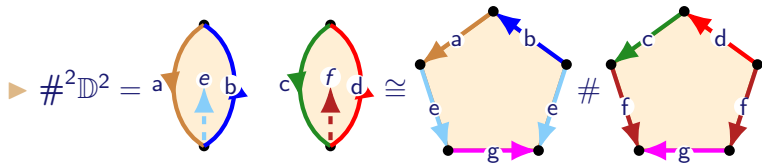
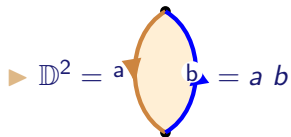
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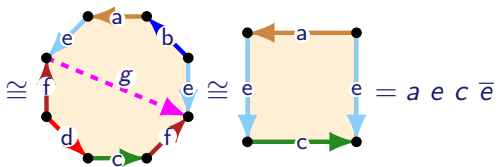
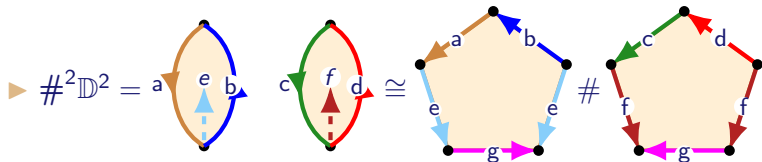
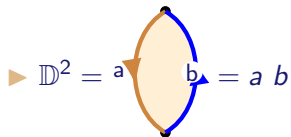
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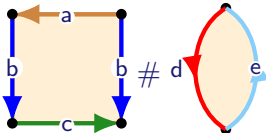
Standard words for surfaces with boundary

▶ $\#^3 \mathbb{D}^2$

Standard words for surfaces with boundary

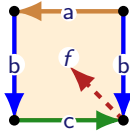
▶ $\#^3 \mathbb{D}^2 = \#^2 \mathbb{D}^2 \# \mathbb{D}^2$

Standard words for surfaces with boundary

► $\#^3 \mathbb{D}^2 = \#^2 \mathbb{D}^2 \# \mathbb{D}^2 \cong$ 

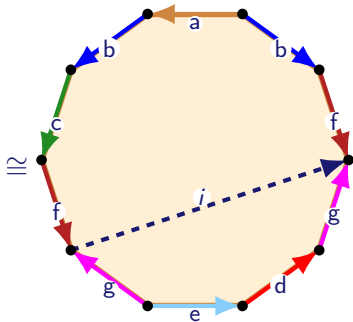
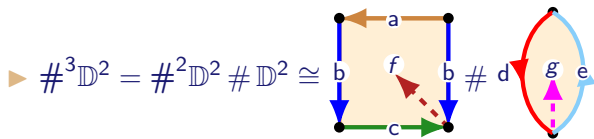
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Standard words for surfaces with boundary

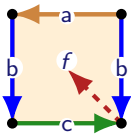
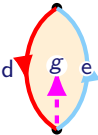
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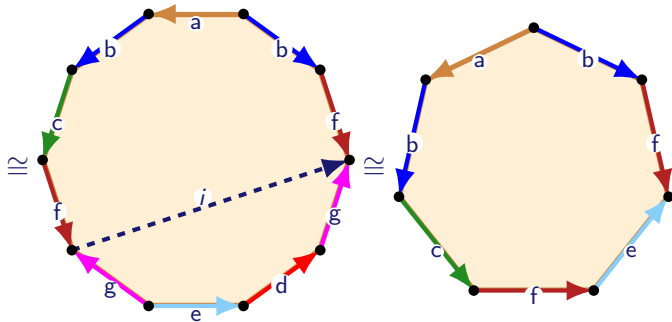
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Standard words for surfaces with boundary

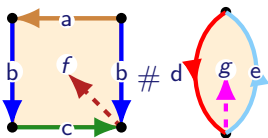


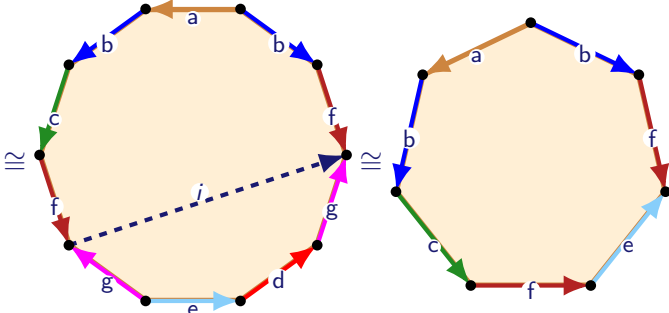
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 $\#$




Standard words for surfaces with boundary

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 \cong

$$= a b c f e \bar{f} \bar{b}$$

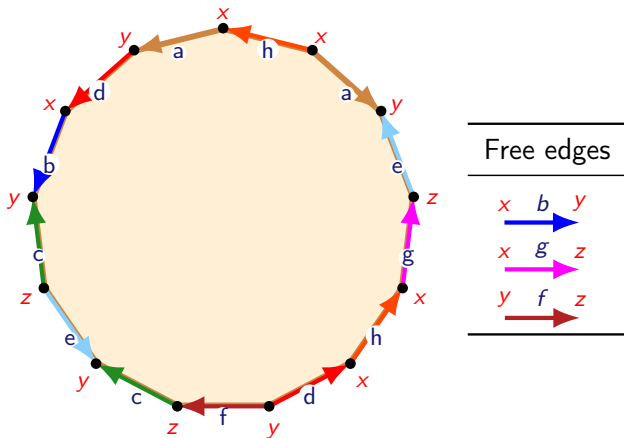
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Words to surfaces

What **standard surface** is given by the word $a d b \bar{c} e \bar{c} \bar{f} d h g e \bar{a} h$?

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$$\implies d = 1 \text{ and } \chi(S) = 3 - 8 + 1 = -4$$

$$\implies S \cong \mathbb{D}^2 \# \#^5 \mathbb{P}^2$$

$$\implies S = a b b c c d d e e f f$$

The vertex-degree equation revisited

When we looked at graphs we proved the **vertex-degree equation**:

$$\sum_{v \in V} \deg(v) = 2|E| \quad \text{for } G = (V, E) \text{ a graph}$$

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The best way to understand this formula is to note that each edge $\{x, y\} \in E$ contributes 2 to both sides of this equation

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- $+2 = 2 \cdot 1$ to the right-hand side for the edge $\{x, w\}$

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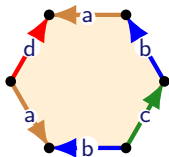
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Answer Yes and no!

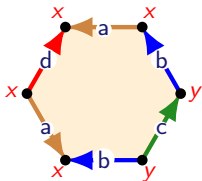
The degree of a vertex

Consider the surface with polygonal decomposition



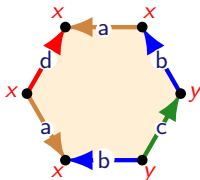
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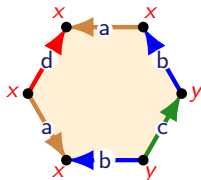
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Using identified vertices and edges + count with multiplicities

The degree of a vertex

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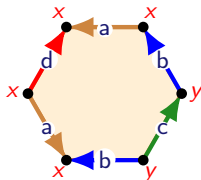


Using identified vertices and edges + count with multiplicities

$$\implies \deg(x) = 5, \deg(y) = 3, \text{ so } \deg(x) + \deg(y) = 8 = 2|E|$$

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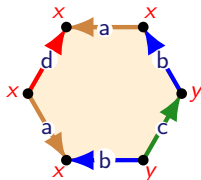
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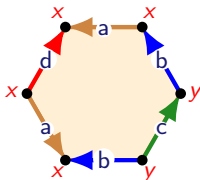
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$$\implies \text{six vertices of degree } 2 \text{ and six edges, so } 12 = 2 \cdot 6$$

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The vertex-degree equation holds using either identified or non-identified edges and vertices because in both cases the **degree** of a vertex is defined to be the **number of incident edges to the vertex**

The surface degree-vertex equation

Proposition

Let $S = (V, E, F)$ be a surface with polygonal decomposition. Then

$$\sum_{v \in V} \deg(v) = 2|E|$$

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Therefore, we have two degree-vertex equations:

- The **graph degree-vertex equation** where we **do not identify** edges and vertices in S
- The **surface degree-vertex equation** where we **do identify** edges and vertices in S

The degree of a face

Let $S = (V, E, F)$ be a surface with polygonal decomposition

Let $f \in F$ be a face of S . The **degree** of f is

$\deg(f)$ = number of edges (count with multiplicities) incident

with f

The degree of a face

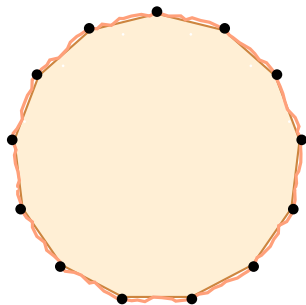
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The degree of a face

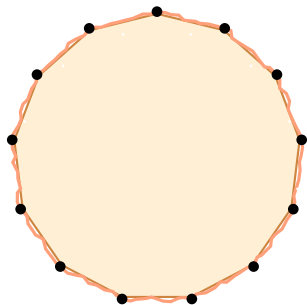
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$$\implies \deg(f) = n$$

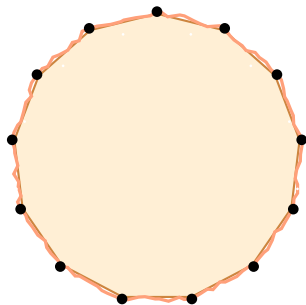
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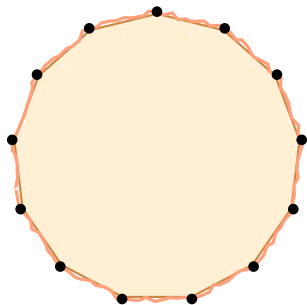
The degree of a face

Let $S = (V, E, F)$ be a surface with polygonal decomposition

Let $f \in F$ be a face of S . The **degree** of f is

$\deg(f)$ = number of edges (count with multiplicities) incident with f

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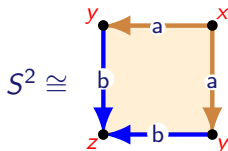
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Question How are $\sum \deg(f)$ and $2|E|$ related?

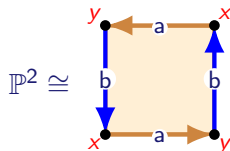
Face degrees of basic surfaces

In all cases $\text{deg}(\text{face}) = 4$ as there are 4 non-identified edges

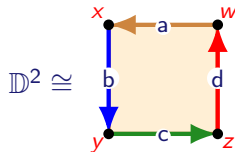
- Sphere



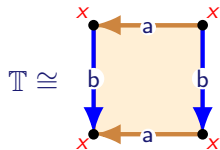
- Projective plane



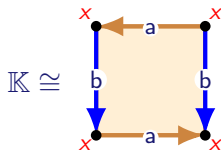
- Disk



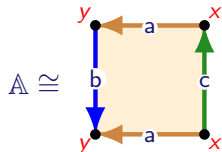
- Torus



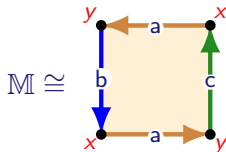
- Klein bottle



- Annulus



- Möbius band



The face-degree equation

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Remark To use this formula we need to know the number of identified edges in the polygonal decomposition

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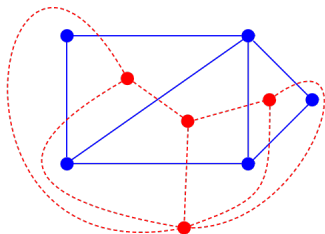
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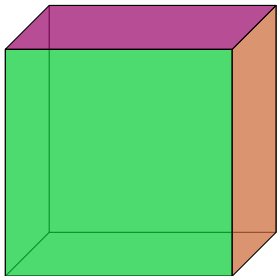
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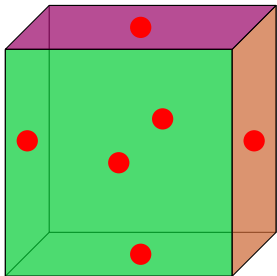
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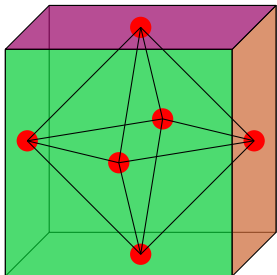
The dual of the cube



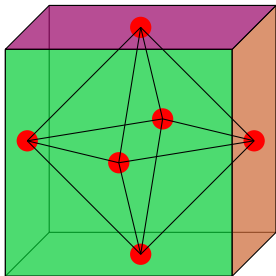
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⇒ the dual surface to the cube is the octahedron

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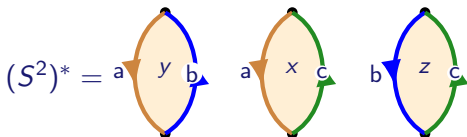
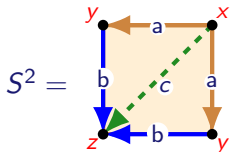
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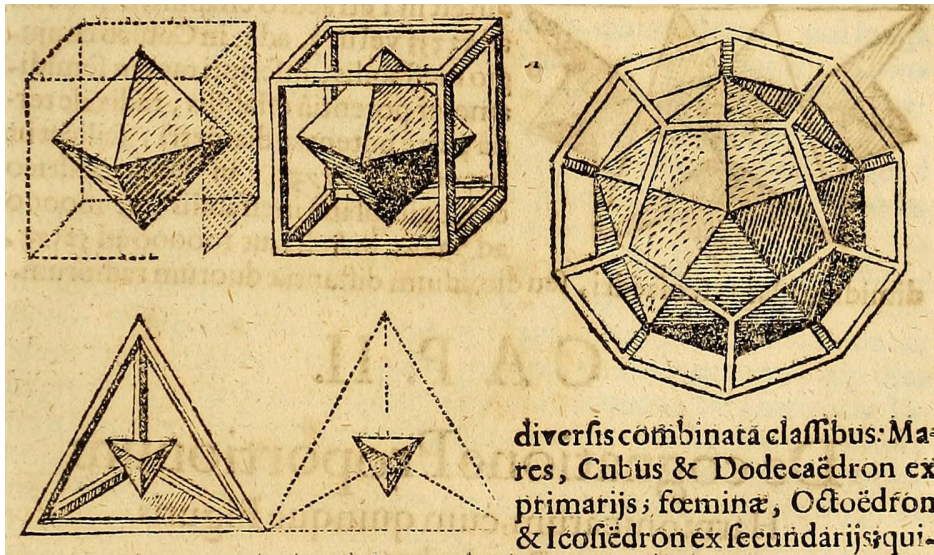
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Example



We will see better examples when we look at Platonic solids

Kepler's Harmonices Mundi



Graphs on surfaces

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- If $e, e' \in E$ then the paths $F(e)$ and $F(e')$ can intersect only at the images of their endpoints

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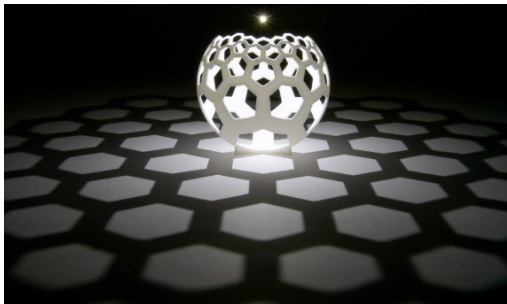
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Proof Stereographic projection! (Move G away from ∞ .)



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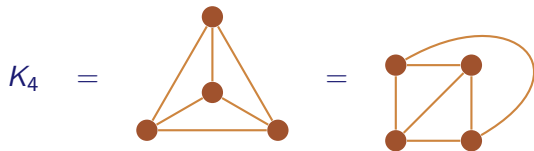
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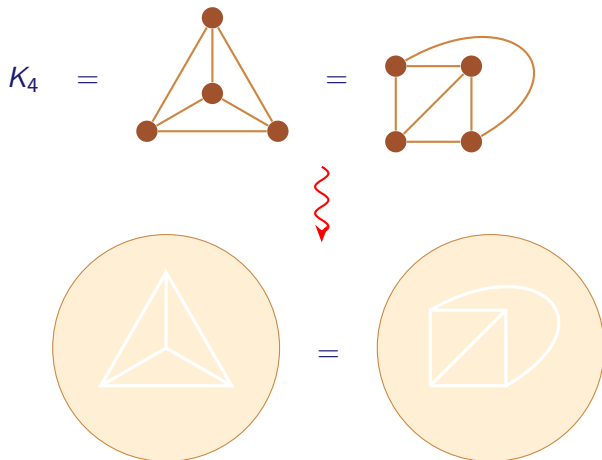
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Remark The argument cheats slightly because we are implicitly assuming that the edges are “nice” curves. This allows us to side-step issues connected with the **Jordan curve theorem**

Planar graphs and Euler characteristic

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Case 1 G is a tree

Combine $|V| - |E| = 1$ (previous lectures) and that there is only one face

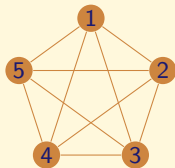
Case 2 G is not a tree

By $\chi(S^2) = 2$ and the previous theorem

Planarity of K_5

Proposition

The complete graph $K_5 =$

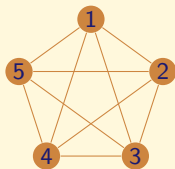


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Proof Assume that K_5 is planar with $|F|$ faces

We have $|V| = 5$ and $|E| = 10$, so $2 = |V| - |E| + |F| \implies |F| = 7$

Let's count the number of faces in this polygonal decomposition differently

- The faces correspond to cycles in K_5
- Every face has at least 3 edges, so by the degree-face equation

$$\implies 2|E| = \sum_{f \in F} \deg(f) \geq 3|F|$$

$$\implies 2|E| = 20 \geq 21 = 3|F| \quad \color{red}{\downarrow \downarrow \downarrow}$$

Hence, the complete graph K_5 is not planar

Planarity of complete graphs

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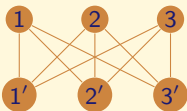
Proof

K_5 sits in K_n for $n \geq 5$, and the previous theorem applies

Planarity of bipartite graphs

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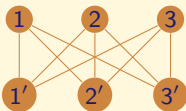


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Planarity of bipartite graphs

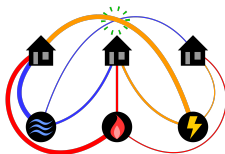
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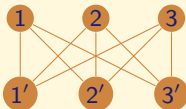
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Proof Tutorials

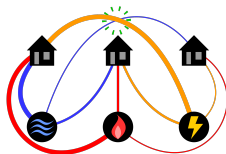


Planarity of bipartite graphs

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Proof Tutorials



Theorem (Kuratowski)

Let G be a graph. Then G is planar if and only if it has no subgraph isomorphic to a **subdivision** of K_5 or $K_{3,3}$

The proof is out of the scope of this unit!






Platonic solids

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



Examples

	Tetrahedron	Cube	Octahedron	Dodecahedron	Icosahedron
					
n	3	4	3	5	3
$ V $	4	8	6	20	12
$ E $	6	12	12	30	30
$ F $	4	6	8	12	20

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n	3	4	3	5	3
$ V $	4	8	6	20	12
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$ F $	4	6	8	12	20

Questions

- Are there any others?

Platonic solids

A **Platonic solid** is a surface that has a polygonal decomposition that is constructed using regular n -gons of the **same shape and size** such that the **same number of polygons meet at every vertex**

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The equations above give:

$$|E| = \left(\frac{1}{p} + \frac{1}{n} - \frac{1}{2} \right)^{-1}, \quad |V| = \frac{2|E|}{p} \quad \text{and} \quad |F| = \frac{2|E|}{n}$$

Classification of Platonic solids

Theorem

The complete list of Platonic solids is:

p	n	$\frac{1}{p} + \frac{1}{n}$	$e = \left(\frac{1}{p} + \frac{1}{n} - \frac{1}{2}\right)^{-1}$	$v = \frac{2e}{p}$	$f = \frac{2e}{n}$	<i>Platonic solid</i>
3	3	$\frac{2}{3}$	6	4	4	<i>Tetrahedron</i>
3	4	$\frac{7}{12}$	12	8	6	<i>Cube</i>
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Proof Since $\frac{1}{p} + \frac{1}{n} > \frac{1}{2}$ and $p, n \geq 3$ we get $n < 6$ since $\frac{1}{3} + \frac{1}{6} = \frac{1}{2}$.
Case-by-case we then get the above values for p, n as the **only possible** values for Platonic solids.

To prove **existence** we need to actually construct them

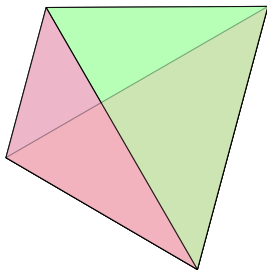
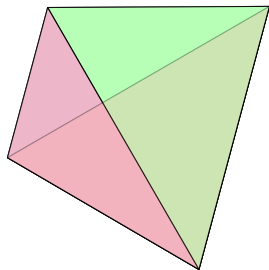
Classification of Platonic solids

Proof Continued Their construction is well-known:

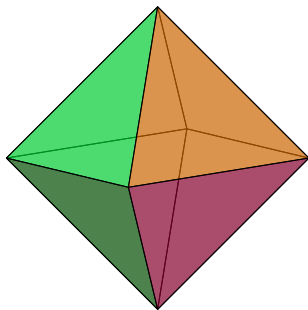
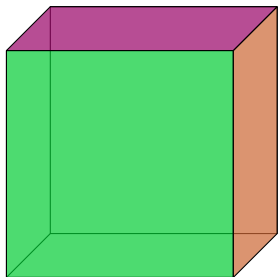


Dual tetrahedron = tetrahedron

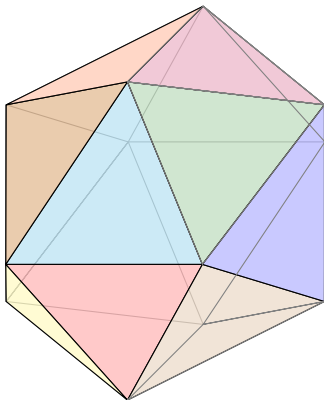
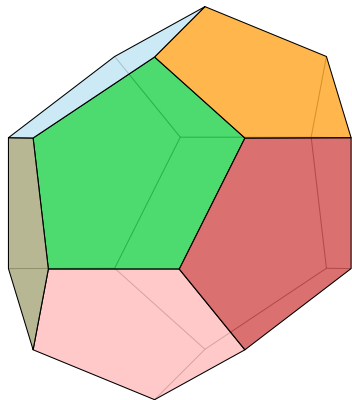
There is a symmetry in the Platonic solids given by $(p, n) \leftrightarrow (n, p)$. This corresponds to taking the dual surface



Cube and octahedron



Dodecahedron and icosahedron



Platonic soccer balls

Here are two dodecahedral decompositions of S^2



Soccer ball

Example A ball is made by gluing together **triangles** and **octagons** so that each octagon is connected to four non-touching triangles. Determine the number of octagons and triangles used

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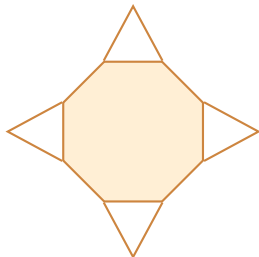
Let there be $|V|$ vertices, $|E|$ edges and $|F|$ faces

Write $|F| = o + t$, where $o = \#$ octagons and $t = \#$ triangles

$$\implies 2 = |V| - |E| + o + t$$

We have:

- vertex-degree equation: $3|V| = 2|E|$
- face-degree equation: $2|E| = 3t + 8o$
- Every octagon meets 4 triangles,
 $\implies 3t = 4o \implies 2|E| = 12o$
 $\implies 2 = o\left(4 - 6 + 1 + \frac{4}{3}\right) = \frac{o}{3}$
 $\implies o = 6$ and $t = 8$
 $\implies |E| = 36$ and $|V| = 24$



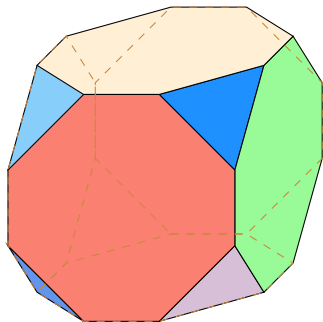
The octacube

As with the Platonic solids, we have only shown that if such a surface exists then there are 6 octagons, 8 triangles, 24 vertices and 36 edges but we have not shown that such a surface exists!

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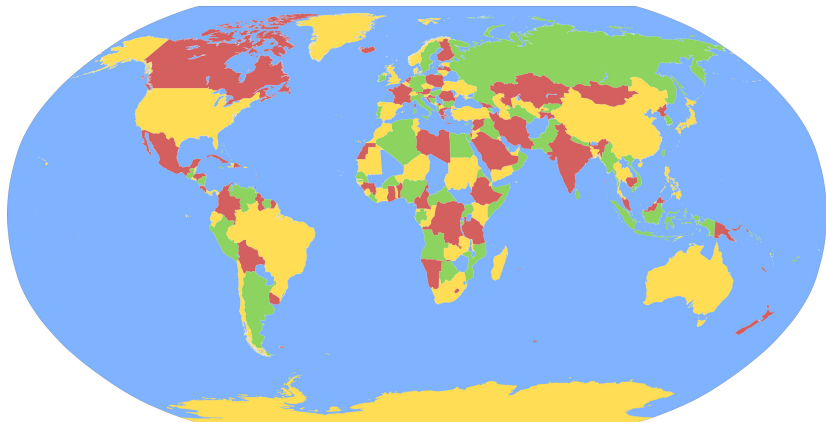
In fact, this surface does exist and it can be constructed by cutting triangular corners off a cube



Coloring maps

Question

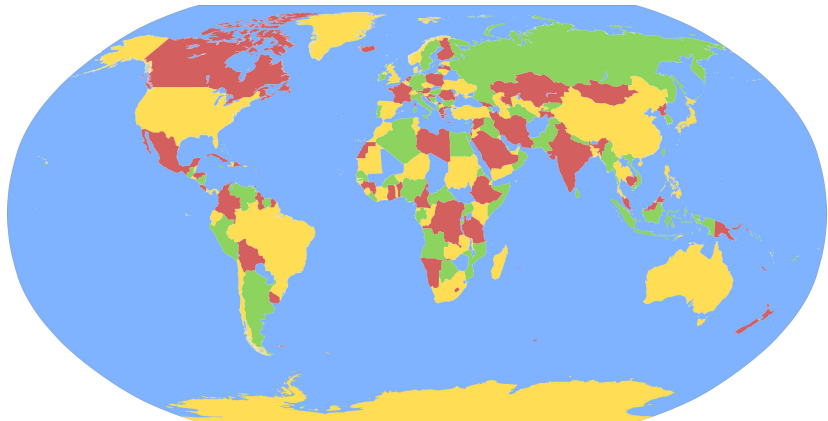
How many different colors do you need to color a map so that adjacent countries have different colors?



Coloring maps

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How many different colors do you need to color a map so that adjacent countries have different colors?



A map is a polygonal decomposition. The answer to this question involves the same ideas we used to understand Platonic solids

Chromatic number of (connected – assumed from now on) surfaces

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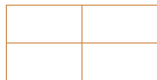
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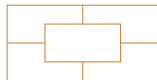
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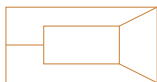
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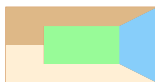
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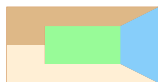
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$$\implies C(\mathbb{D}^2) \geq 4$$

For maps of the world we are most interested in $C(\mathbb{D}^2) = C(S^2)$

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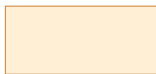
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These assumptions are purely for convenience because, in each case, we can colour these maps using the same number of colours

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The basic idea is to use the Euler characteristic and the degree-vertex and degree-face equations to understand colourings

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Remark For a Platonic solid that is made from n -gons with p polygons meeting at each vertex we have $\partial_V = p$ and $\partial_F = n$

Bounding the face degree

Lemma

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Maps on sphere and projective planes

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- 2 If the average face degree $\partial_F < 6$ then there must be at least one face f with $\deg(f) \leq 5$
This observation will be important when we prove the **Five color theorem** (not quite the four color theorem)