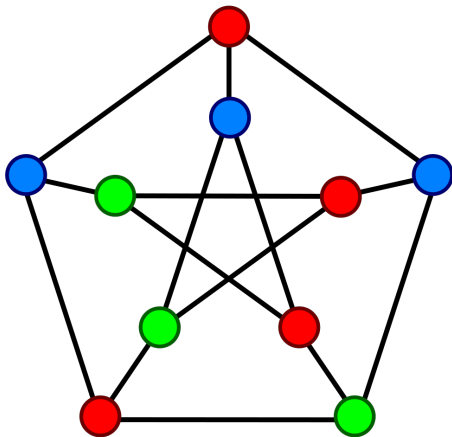


What is...the chromatic polynomial?

Or: Polynomials and colors

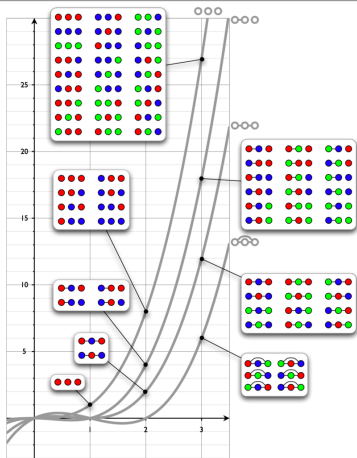
Counting colorings

3-coloring:



-
- ▶ Coloring = vertices get colors such that two adjacent vertices have different colors
 - ▶ Let $P_G(k)$ = number of k -colorings of G (colorings using k colors)
 - ▶ Question How are the $P_G(k)$ related?

3 vertex graph



edgeless: x^3

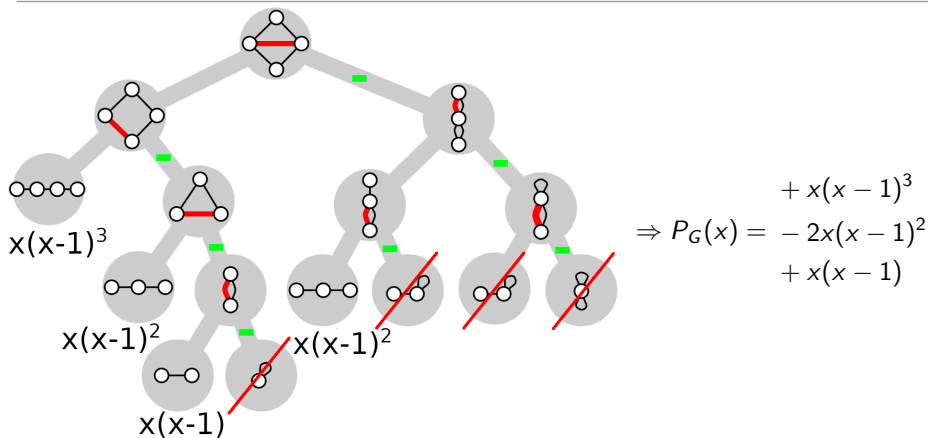
one edge: $x^2(x - 1)$

line: $x(x - 1)^2$

circle: $x(x - 1)(x - 2)$

- ▶ For $G = \bullet \bullet \bullet$ we have a polynomial giving $P_G(k)$, namely $P_G(x) = x^3$
- ▶ For other small graphs one checks that the same works
- ▶ Question Is there a polynomial counting colorings?

Deletion-contraction



- ▶ Here is an **algorithm** to compute $P_G(x)$
- ▶ **Starting condition** $P_{tree}(x) = x(x-1)^{\#vertices-1}$ and $P_{loop}(x) = 0$
- ▶ Then use **deletion-contraction**: $P_G(x) = P_{G \setminus e}(x) - P_{G/e}(x)$

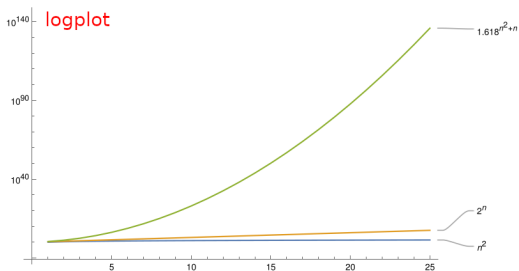
For completeness: A formal statement

There exists a polynomial $P_G(x)$ associated to a graph such that:

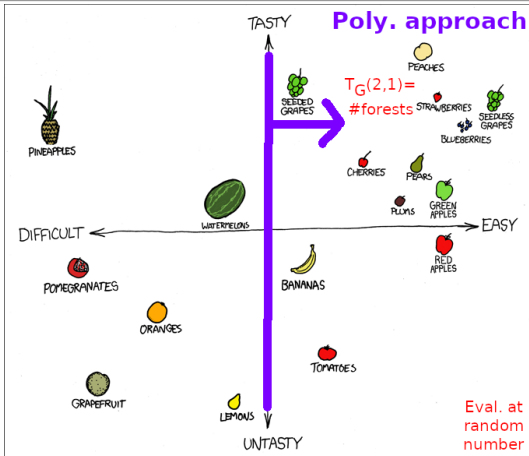
$$P_G(k) = \#k\text{-colorings}$$

- ▶ The polynomial is called chromatic polynomial
- ▶ The polynomial can be computed by deletion-contraction
- ▶ However, the runtime is quite bad: $\approx \phi^{\#vertices + \#edges}$; $\phi = \text{golden ratio}$

$\#edges \approx \#vertices^2$:



Whatever “easy” means



- ▶ Recall that graph polynomials are for **easy problems**
- ▶ The runtime for $P_G(x)$ is **horrible**, so how can that be easy?
- ▶ This is easy in the sense that we get **all** colorings at once

Thank you for your attention!

I hope that was of some help.