

What is...the complexity of the Tutte polynomial?

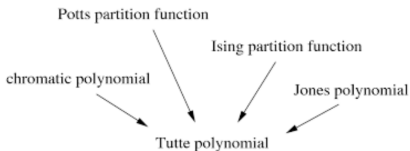
Or: Easy or difficult?

Computing the Tutte polynomial $T_G(x, y)$

For completeness: A formal statement

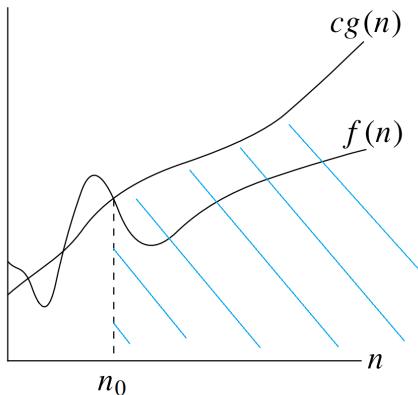
There exists a polynomial $T_G(x, y)$ associated to a graph such that:

- ▶ $T_G(2, 1) = \#$ forests
 - ▶ $T_G(1, 1) = \#$ spanning forests
 - ▶ $T_G(1, 2) = \#$ spanning subgraphs
 - ▶ More...
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- ▶ The polynomial is called Tutte polynomial
 - ▶ Also we have the specialization “chromatic(x) = Tutte(x,0)”, and more



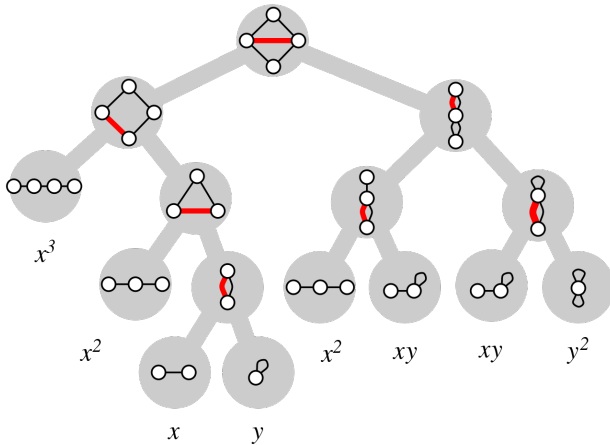
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- ▶ $T_G(x, y)$ counts many things, so it would be good to compute it efficiently
 - ▶ Question How difficult is it to compute $T_G(x, y)$?
 - ▶ Question How difficult is it to compute $T_G(a, b)$ for $(a, b) \in \mathbb{C}^2$?

Landau–Bachmann notation



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- ▶ We say $f \in O(g)$ if $f(n) \leq cg(n)$, $c = \text{constant}$, from some point onward
 - ▶ Example $10000n \in O(n^2)$
 - ▶ We use this to analyze **worst-case runtime** for algorithms

The computation via recursion

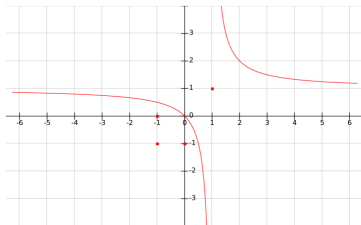


- ▶ Recall the deletion-contraction way to compute $T_G(x, y)$
- ▶ This looks like exponential growth
- ▶ **Guess** Computing $T_G(x, y)$ is probably difficult e.g. Tutte $\in O(2^{\#edges})$

For completeness: A formal statement

The computation of $T_G(a, b)$ is...

- ▶ ...in $O(\text{polynomial})$ for $(a - 1)(b - 1) = 1$ Easy



- ▶ ...in $O(\text{polynomial})$ for $(j = \exp(2\pi i/3))$
 $(a, b) \in \{(1, 1), (-1, -1), (0, -1), (-1, 0), (i, -i), (-i, i), (j, j^2), (j^2, j)\}$ Easy

- ▶ ...#P hard otherwise Hard

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- ▶ #P hard \approx Tutte $\in O(2^{\#\text{edges}})$ but the precise runtime is unknown
 - ▶ Note the huge difference between general and specific points

Difficult in general, but...

Graph class	#P-hard	subexponential	FPT	P
All graphs	$\mathbb{C}^2 - H$	H	H	H
planar	$\mathbb{C}^2 - H_2$	H_2	H_2	H_2
bipartite planar	$\mathbb{C}^2 - H_{b-p}$	H_{b-p}	H_{b-p}	H_{b-p}
$TW(k)$	\emptyset	\mathbb{C}^2	\mathbb{C}^2	H
$CW(k)$	\emptyset	\mathbb{C}^2	H	H

H = hyperbola from the previous slide

$TW(k)$ = tree width at most k

$CW(k)$ = clique width at most k

► Computing $T_G(x, y)$ in general is difficult

► Computing $T_G(x, y)$ in special cases is not so bad

Thank you for your attention!

I hope that was of some help.