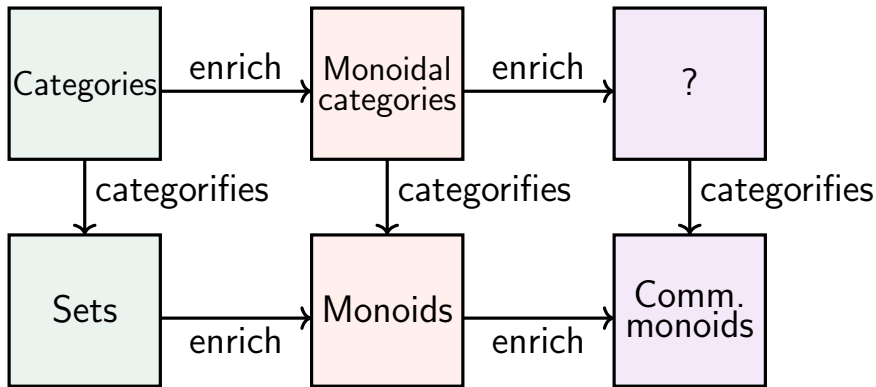


What are...braided categories?

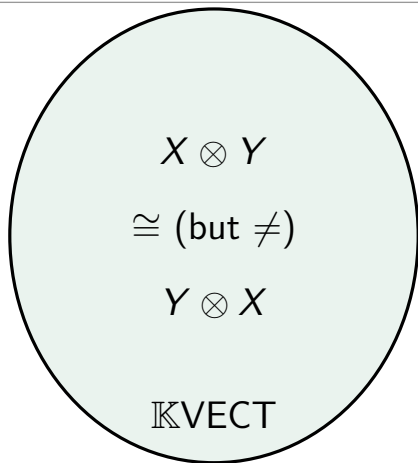
Or: Braids and categories

Filling in more question marks



- ▶ Monoidal categories categorify monoids
- ▶ Commutative monoids are enriched monoids
- ▶ What completes the rectangle?

Vector spaces again (what else?)



-
- ▶ In $\mathbb{K}\text{VECT}$ we have $X \otimes Y \cong Y \otimes X$
 - ▶ The isomorphism is the swap map $\beta_{X,Y}: x \otimes y \mapsto y \otimes x$
 - ▶ This looks like **commutativity**

String diagram and braids

$$\beta_{X,Y} \iff \begin{array}{c} Y \quad X \\ \diagdown \quad / \\ \diagup \quad \diagdown \\ X \quad Y \end{array}, \quad \beta_{X,Y}^{-1} \iff \begin{array}{c} Y \quad X \\ \diagup \quad \diagdown \\ \diagdown \quad / \\ X \quad Y \end{array}$$

actually for $\mathbb{K}\text{VECT}$: $\beta_{X,Y} = \beta_{X,Y}^{-1} \iff \begin{array}{c} Y \quad X \\ \diagdown \quad / \\ \diagup \quad \diagdown \\ X \quad Y \end{array} = \begin{array}{c} Y \quad X \\ \diagup \quad \diagdown \\ \diagdown \quad / \\ X \quad Y \end{array} = \begin{array}{c} Y \quad X \\ \diagdown \quad / \\ \diagdown \quad / \\ X \quad Y \end{array}$

$$\begin{array}{c} X \quad Y \\ \diagdown \quad / \\ \diagup \quad \diagdown \\ X \quad Y \end{array} = \begin{array}{c} X \quad Y \\ | \quad | \\ X \quad Y \end{array}, \quad \begin{array}{c} Y \quad X \\ \diagdown \quad / \\ \diagup \quad \diagdown \\ Y \quad X \end{array} = \begin{array}{c} Y \quad X \\ | \quad | \\ Y \quad X \end{array}, \quad \begin{array}{c} Z \quad Y \quad X \\ \diagdown \quad / \quad \diagdown \\ \diagup \quad \diagdown \quad / \\ X \quad Y \quad Z \end{array} = \begin{array}{c} Z \quad Y \quad X \\ \diagup \quad \diagdown \quad / \\ \diagdown \quad / \quad \diagdown \\ X \quad Y \quad Z \end{array}$$

- ▶ Denote the swap map by a crossing
- ▶ Categorically commutativity looks like a braid group action

For completeness: A formal definition

A monoidal category C is braided (strict) (extra data!) if:

- There exists a collection of natural isomorphism

$$\beta_{X,Y}: XY \xrightarrow{\cong} YX, \text{ think } \beta_{X,Y} \leftrightarrow \begin{array}{c} Y \ X \\ \swarrow \searrow \\ X \ Y \end{array}, \quad \beta_{X,Y}^{-1} \leftrightarrow \begin{array}{c} Y \ X \\ \swarrow \searrow \\ X \ Y \end{array}$$

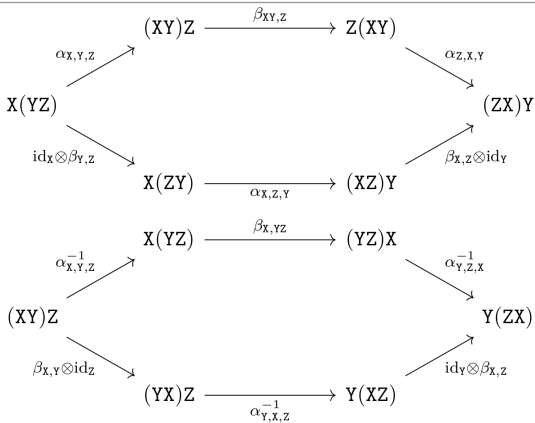
- We have naturality and sliding type relations (these imply the braid relations):

$$\begin{array}{c} Z \ X \ Y \\ \swarrow \searrow \\ X \ Y \ Z \end{array} = \begin{array}{c} Z \ \ \ XY \\ \swarrow \searrow \\ X \ Y \ Z \end{array}, \quad \begin{array}{c} Y \ Z \ X \\ \swarrow \searrow \\ X \ Y \ Z \end{array} = \begin{array}{c} YZ \ \ X \\ \swarrow \searrow \\ X \ \ \ YZ \end{array}, \quad \begin{array}{c} A \ \ Y \\ \swarrow \searrow \\ \boxed{f} \ \ \boxed{g} \\ | \ \ \ | \\ X \ \ \ Z \end{array} = \begin{array}{c} A \ \ \ Y \\ \swarrow \searrow \\ \boxed{g} \ \ \boxed{f} \\ | \ \ \ | \\ X \ \ \ Z \end{array}$$

A category is symmetric (strict) if it is braided with

$$\beta_{X,Y} = \beta_{X,Y}^{-1}, \text{ think } \beta_{X,Y} = \beta_{X,Y}^{-1} \leftrightarrow \begin{array}{c} Y \ X \\ \swarrow \searrow \\ X \ Y \end{array} = \begin{array}{c} Y \ X \\ \swarrow \searrow \\ X \ Y \end{array} = \begin{array}{c} Y \ X \\ \swarrow \searrow \\ X \ Y \end{array}$$

Strict vs. not strict



- ▶ The non strict definition of braided involves several big commutative diagrams
- ▶ As usual, there is a **strictification** result

Every braided category is braided equivalent to a strict braided category

so I ignored the difference

Thank you for your attention!

I hope that was of some help.