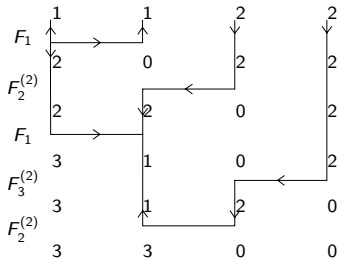


\mathfrak{sl}_3 -web bases, categorification and link invariants

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1 Categorification

- What is categorification?

2 Webs and representations of $U_q(\mathfrak{sl}_3)$

- \mathfrak{sl}_3 -webs
- Intermediate crystals
- q -skew Howe duality and growing of webs
- Knot polynomials and webs

3 The Categorification

- An algebra of foams
- Growing of foams
- Foamation
- Knot homologies and foams

What is categorification?

Forced to reduce this presentation to one sentence, the author would choose:

Interesting integers are shadows of richer structures in categories.

The basic idea can be seen as follows. Take a “set-based” structure S and try to find a “category-based” structure \mathcal{C} such that S is just a shadow of \mathcal{C} .

Categorification, which can be seen as “remembering” or “inventing” information, comes with an “inverse” process called decategorification, which is more like “forgetting” or “identifying”.

Note that decategorification should be easy.

Exempli gratia

Examples of the pair categorification/decategorification are:

Bettinnumbers of a manifold M	$\begin{array}{c} \xrightarrow{\text{categorify}} \\ \xleftarrow{\text{decat}=\text{rank}(\cdot)} \end{array}$	Homology groups
Polynomials in $\mathbb{Z}[q, q^{-1}]$	$\begin{array}{c} \xrightarrow{\text{categorify}} \\ \xleftarrow{\text{decat}=\chi_{\text{gr}}(\cdot)} \end{array}$	complexes of gr.VS
The integers \mathbb{Z}	$\begin{array}{c} \xrightarrow{\text{categorify}} \\ \xleftarrow{\text{decat}=\mathcal{K}_0(\cdot)} \end{array}$	K – vector spaces
An A – module	$\begin{array}{c} \xrightarrow{\text{categorify}} \\ \xleftarrow{\text{decat}=\mathcal{K}_0^{\oplus}(\cdot) \otimes_{\mathbb{Z}} A} \end{array}$	additive category

Usually the **categorified world** is much more **interesting**.

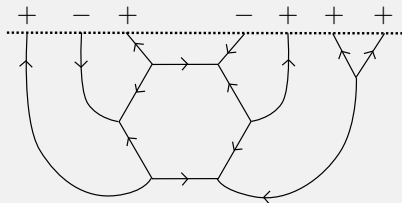
Today decategorification = Grothendieck group!

Kuperberg's \mathfrak{sl}_3 -webs

Definition(Kuperberg)

A \mathfrak{sl}_3 -web w is an **oriented trivalent graph**, such that all vertices are either sinks or sources. The boundary ∂w of w is a **sign string** $S = (s_1, \dots, s_n)$ under the convention $s_i = +$ iff the orientation is pointing in and $s_i = -$ iff the orientation is pointing out (we also need 0, 3 later - but they are **not** drawn).

Example



Definition(Kuperberg)

The $\mathbb{C}(q)$ -web space W_S for a given sign string $S = (\pm, \dots, \pm)$ is generated by $\{w \mid \partial w = S\}$, where w is a web, subject to the relations

$$\begin{array}{l}
 \text{circle} = [3] \\
 \text{line with loop} = [2] \text{ line} \\
 \text{square with arrows} = \left. \begin{array}{l} \text{left brace} \\ \text{right brace} \end{array} \right\} + \text{crossing}
 \end{array}$$

Here $[a] = \frac{q^a - q^{-a}}{q - q^{-1}} = q^{a-1} + q^{a-3} + \dots + q^{-(a-1)}$ is the **quantum integer**.

A connection to knot theory

Let L_D be a link projection. Assign to it a **polynomial** $P_3(L_D)$ (it is in $\mathbb{Z}[q, q^{-1}]$) by local and inductive rules as follows.

- $P_3(\text{crossing}) = q^2 P_3(\text{positive}) - q^3 P_3(\text{negative})$ (recursion rule 1).
- $P_3(\text{crossing}) = q^{-2} P_3(\text{positive}) - q^{-3} P_3(\text{negative})$ (recursion rule 2).
- The Kuperberg relations.

Theorem (Murakami, Ohtsuki and Yamada)

The polynomial $P_3(\cdot)$ is uniquely determined by the rule and a link **invariant**. Moreover, it **agrees** with the so-called HOMFLY-PT polynomial under a certain substitution of variables and normalization.

For example

$$\begin{aligned} P_3 \left(\text{link with two crossings} \right) &= q^2 P_3 \left(\text{link with two crossings} \right) - q^3 P_3 \left(\text{link with two crossings} \right) \\ &= q^2 [3]^2 - q^3 [2][3] = [3]. \end{aligned}$$

Representation theory of $\mathbf{U}_q(\mathfrak{sl}_3)$

A sign string $S = (s_1, \dots, s_n)$ corresponds to tensors

$$V_S = V_{s_1} \otimes \cdots \otimes V_{s_n},$$

where V_+ is the fundamental $\mathbf{U}_q(\mathfrak{sl}_3)$ -representation and V_- is its dual, and webs correspond to **intertwiners**.

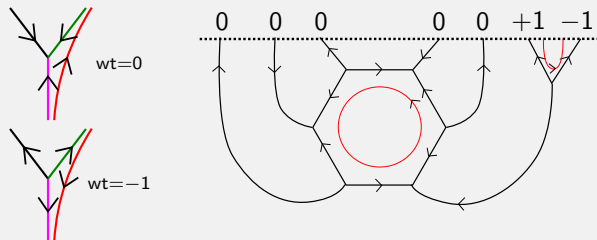
Theorem(Kuperberg)

$$W_S \cong \text{Hom}_{\mathbf{U}_q(\mathfrak{sl}_3)}(\mathbb{C}(q), V_S) \cong \text{Inv}_{\mathbf{U}_q(\mathfrak{sl}_3)}(V_S)$$

In fact, the so-called spider category of all webs modulo the Kuperberg relations is **equivalent** to the representation category of $\mathbf{U}_q(\mathfrak{sl}_3)$.

As a matter of fact, the \mathfrak{sl}_3 -webs without internal circles, digons and squares form a **basis** B_S , called **web basis**, of W_S !

Example



Webs can be coloured with flow lines. At the boundary, the flow lines can be represented by a state string J . By convention, at the i -th boundary edge, we set $j_i = \pm 1$ if the flow line is oriented downward/upward and $j_i = 0$, if there is no flow line. So $J = (0, 0, 0, 0, 0, +1, -1)$ in the example.

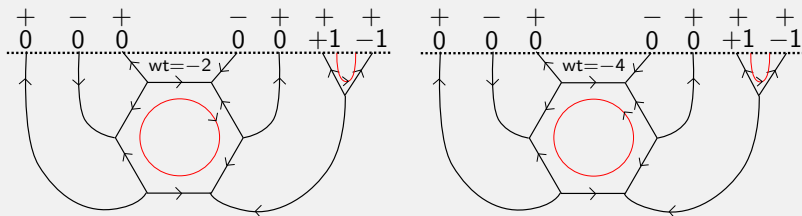
Given a web with a flow w_f , attribute a weight to each trivalent vertex and each arc in w_f and take the sum. The weight of the example is -4 .

Representation theory of $U_q(\mathfrak{sl}_3)$

Theorem (Khovanov-Kuperberg)

Pairs of sign S and a state strings J correspond to the coefficients of the web basis relative to **tensors of the standard basis** $\{e_{\pm}^{-1}, e_{\pm}^0, e_{\pm}^{+1}\}$ of V_{\pm} .

Example



$$w_S = \dots + (q^{-2} + q^{-4})(e_+^0 \otimes e_-^0 \otimes e_+^0 \otimes e_-^0 \otimes e_+^0 \otimes e_+^1 \otimes e_+^{-1}) \pm \dots$$

What kind of basis is B_S ?

Theorem(Khovanov-Kuperberg)

Given (S, J) , we have (with $v = -q^{-1}$ and $e_S^J = e_{s_1}^{j_1} \otimes \cdots \otimes e_{s_n}^{j_n}$)

$$w_S^J = e_S^J + \sum_{J' < J} c(S, J, J') e_S^{J'} \quad \text{for } c(S, J, J') \in \mathbb{N}[v, v^{-1}].$$

In general we have $B_S \neq \text{dcan}(W_S)$, but the web basis is **bar-invariant**.

Theorem(MPT)

We proved, by categorification, that the change-of-basis matrix from Kuperberg's web basis B_S to the dual canonical basis $\text{dcan}(W_S)$ is **unitriangular**.

Question: The web basis B_S is a somehow "special" basis of W_S . But it is **not** the dual canonical. So what kind of basis is it?

The quantum algebra $\mathbf{U}_q(\mathfrak{sl}_d)$

Definition

For $d \in \mathbb{N}_{>1}$ the **quantum special linear algebra** $\mathbf{U}_q(\mathfrak{sl}_d)$ is the associative, unital $\mathbb{C}(q)$ -algebra generated by $K_i^{\pm 1}$ and E_i and F_i , for $i = 1, \dots, d-1$, subject to some relations (that we do not need today).

Definition (Beilinson-Lusztig-MacPherson)

For each $\lambda \in \mathbb{Z}^{d-1}$ adjoin an **idempotent** 1_λ (**think**: projection to the λ -weight space!) to $\mathbf{U}_q(\mathfrak{sl}_d)$ and add some relations, e.g.

$$1_\lambda 1_\mu = \delta_{\lambda, \mu} 1_\lambda \quad \text{and} \quad K_{\pm i} 1_\lambda = q^{\pm \lambda_i} 1_\lambda \quad (\text{no } K\text{'s anymore!}).$$

The **idempotent quantum special linear algebra** is defined by

$$\dot{\mathbf{U}}(\mathfrak{sl}_d) = \bigoplus_{\lambda, \mu \in \mathbb{Z}^{d-1}} 1_\lambda \mathbf{U}_q(\mathfrak{sl}_d) 1_\mu.$$

Intermediate crystals

Let $d = 3\ell$ and let V_Λ be the irreducible $\dot{\mathbf{U}}(\mathfrak{sl}_d)$ -module of highest weight $\Lambda = (3^\ell)$. Kashiwara-Lusztig's **lower global crystal** (or **canonical**) basis $\text{can}(V_\Lambda) = \{b_T \mid T \in \text{Std}(3^\ell)\}$ has **nice** properties, but is in **very hard** to find.

Leclerc and Toffin have defined an **intermediate** crystal basis B_Λ of V_Λ by an explicit algorithm that can be used to compute $\text{can}(V_\Lambda)$ inductively, i.e. B_Λ has **some nice** properties, but is still **trackable** enough to be written down.

Example (with $\ell = 3$)

$$T = \begin{array}{|c|c|c|} \hline 1 & 2 & 4 \\ \hline 2 & 3 & 5 \\ \hline 4 & 6 & 6 \\ \hline \end{array} \rightsquigarrow \text{LT}(T) = F_1 F_2 F_3^{(2)} F_2 F_1 F_4 F_3 F_2 F_5^{(2)} F_4^{(2)} F_3^{(2)} v_\Lambda.$$

Sitting in-between $\{b_T\}$ and $\{x_{T'}\}$

The intermediate crystal basis sits “in-between” the canonical $\text{can}(V_\Lambda)$ and the tensor basis $\{x_{T'} \in \Lambda_q^\ell(\mathbb{C}_q^d)^{\otimes 3} \supset V_\Lambda \mid T' \in \text{Col}(3^\ell)\}$.

Theorem(Leclerc-Toffin)

We have (for $T' \in \text{Col}(3^\ell)$, $T'' \in \text{Std}(3^\ell)$)

$$\text{LT}(T) = x_T + \sum_{T' \prec T} \alpha_{T'}(v)x_{T'} \quad \text{and} \quad b_T = \text{LT}(T) + \sum_{T'' \prec T} \beta_{T''T}(v)\text{LT}(T'')$$

with certain $\alpha_{T'}(v) \in \mathbb{N}[v, v^{-1}]$ and $\beta_{T''T}(v) \in \mathbb{Z}[v, v^{-1}]$ (with $v = -q^{-1}$).
Moreover, the intermediate crystal basis is **bar-invariant**.

We have seen this **before!**

An instance of q -skew Howe duality

The commuting actions of $\dot{\mathbf{U}}(\mathfrak{sl}_d)$ and $\dot{\mathbf{U}}(\mathfrak{sl}_3)$ on

$$\Lambda_q^\bullet(\mathbb{C}_q^d)^{\otimes 3} \cong \Lambda_q^\bullet(\mathbb{C}^d \otimes \mathbb{C}^3) \cong \Lambda_q^\bullet(\mathbb{C}_q^3)^{\otimes d}$$

introduce an $\dot{\mathbf{U}}(\mathfrak{sl}_d)$ -action on $\Lambda_q^\bullet(\mathbb{C}_q^3)^{\otimes d}$ and an $\dot{\mathbf{U}}(\mathfrak{sl}_3)$ -action on $\Lambda_q^\bullet(\mathbb{C}_q^d)^{\otimes 3}$. Here

$$\Lambda_q^\bullet(\mathbb{C}_q^l) = \bigoplus_{k=1}^l \Lambda_q^k(\mathbb{C}_q^l)$$

and all the $\Lambda_q^k(\mathbb{C}_q^l)$ are irreducible $\dot{\mathbf{U}}(\mathfrak{sl}_l)$ -representations.

The \mathfrak{sl}_3 -webs form a $\dot{\mathbf{U}}(\mathfrak{sl}_d)$ -module

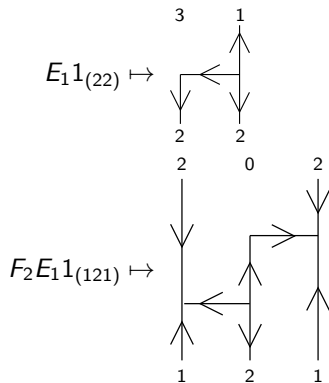
We **defined** an action ϕ of $\dot{\mathbf{U}}(\mathfrak{sl}_d)$ on $W_{(3^\ell)} = \bigoplus_{S \in \Lambda(n, n)_3} W_S$ by

$$\begin{array}{c}
 1_\lambda \mapsto \begin{array}{c} | & | & \dots & | \\ \lambda_1 & \lambda_2 & & \lambda_d \end{array} \\
 \\
 E_i 1_\lambda, F_i 1_\lambda \mapsto \begin{array}{c} \lambda_i \pm 1 \quad \lambda_{i+1} \mp 1 \\ | & | & | & | & | & | \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \lambda_1 & \lambda_{i-1} & \lambda_i & \lambda_{i+1} & \lambda_{i+2} & \lambda_d \end{array}
 \end{array}$$

We use the convention that vertical edges labeled 1 are oriented upwards, vertical edges labeled 2 are oriented downwards and edges labeled 0 or 3 are erased.

Think: 0, 3 indicates the trivial, 1 the V_+ and 2 the V_- -representation.

Exempli gratia



An intermediate crystal basis

Proposition(MT)

The Kuperberg web basis B_S is Leclerc-Toffin's intermediate crystal basis under q -skew Howe duality, i.e.

$$\text{LT}(T) = \{F_{i_s}^{(r_s)} \cdots F_{i_1}^{(r_1)} v_{3^\ell} \mid T \in \text{Std}(3^\ell)\} \xrightarrow{\text{sHd}} w_S^J.$$

(No K 's and E 's anymore!)

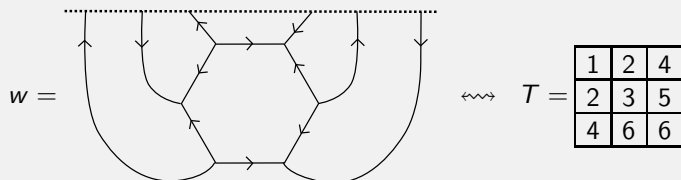
Corollary

The change-of-basis matrix from Kuperberg's web basis B_S to the dual canonical basis $\text{dcan}(W_S)$ is **unitriangular**, it is bar-invariant and

$$w_S^J = e_S^J + \sum_{J' < J} c(S, J, J') e_S^{J'} \quad \text{for } c(S, J, J') \in \mathbb{N}[v, v^{-1}].$$

We also defined a **growth algorithm for flows** - but we do not need it today.

Example

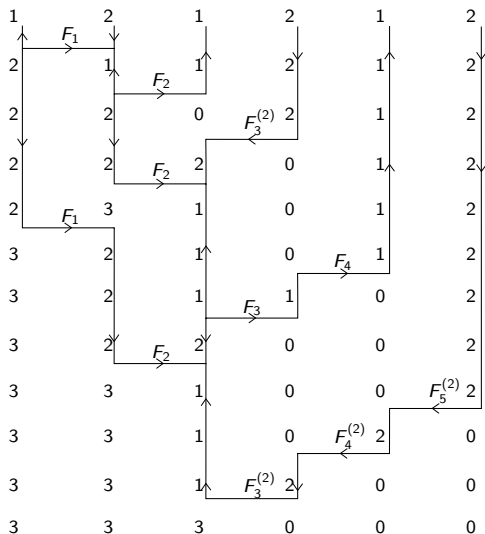


From T we obtain the string

$$\text{LT}(T) = F_1 F_2 F_3^{(2)} F_2 F_1 F_4 F_3 F_2 F_5^{(2)} F_4^{(2)} F_3^{(2)}.$$

Exempli gratia

$$LT(T)v_{3^3} = F_1 F_2 F_3^{(2)} F_2 F_1 F_4 F_3 F_2 F_5^{(2)} F_4^{(2)} F_3^{(2)} v_{3^3}$$



Crossings as $F_j^{(i)}$

Define the following **braiding** operators B_j (and **similar ones** for $\searrow \swarrow$).

$$\swarrow \nearrow : (q^2 F_j F_{j+1} - q^3 F_{j+1} F_j) v_{110} \quad \searrow \swarrow : (q^{-2} F_{j+1}^{(2)} F_j^{(2)} - q^{-3} F_{j+1} F_j^{(2)} F_{j+1}) v_{210}$$

$$\searrow \swarrow : (q^2 F_{j+1} F_j - q^3 F_j F_{j+1}) v_{332} \quad \swarrow \searrow : (q^{-2} F_{j+1} F_j - q^{-3} F_j F_{j+1}) v_{120}$$

Example

$$\swarrow \searrow : (q^{-2} F_{j+1} F_j - q^{-3} F_j F_{j+1}) v_{120}$$

gives (for $j = 1$)

$$\begin{array}{ccc}
 \begin{array}{c} 0 \\ \swarrow \searrow \\ \nearrow \swarrow \\ 1 \end{array} & \begin{array}{c} 2 \\ \searrow \swarrow \\ \nearrow \swarrow \\ 3 \end{array} & \begin{array}{c} 1 \\ \swarrow \searrow \\ \nearrow \swarrow \\ 0 \end{array} \\
 q^{-2} & & -q^{-3} \\
 \begin{array}{c} 0 \\ \swarrow \searrow \\ \nearrow \swarrow \\ 1 \end{array} & \begin{array}{c} 2 \\ \searrow \swarrow \\ \nearrow \swarrow \\ 2 \end{array} & \begin{array}{c} 1 \\ \swarrow \searrow \\ \nearrow \swarrow \\ 0 \end{array}
 \end{array}$$

Proposition(T)

Let L_D be a diagram of a link. Then (under q -skew Howe duality):

$$\text{LT}(L_D)v_h = \prod_k \tilde{F}_{j_k}^{(i_k)} v_h, \quad \text{with } \tilde{F} = F, \tilde{F} = B$$

for some highest weight vector v_h . That is, L_D can be realized **combinatorial** as sums of certain tableaux.

Observation(T)

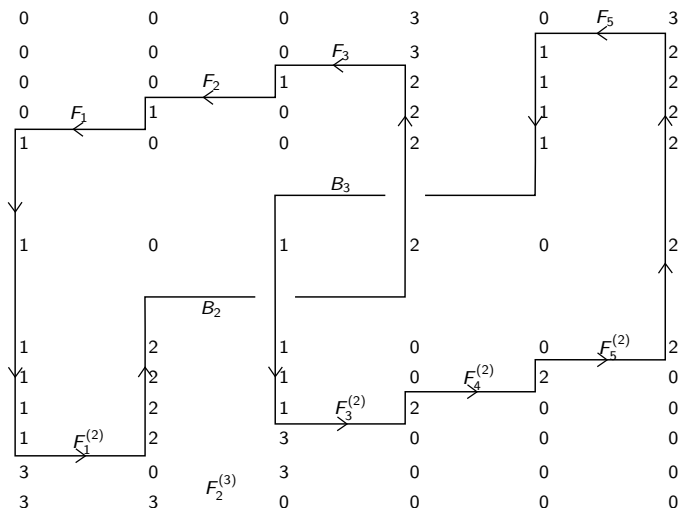
Since the above construction agrees (up to some normalization/shifts) with the construction of the \mathfrak{sl}_3 -link polynomial (MOY-calculus), the combinatoric of tableaux can be used to calculate these invariants (note: works for **all** $n > 1$).

Wish(T)

There should be a way to get the **colored** \mathfrak{sl}_3 -link polynomial from this approach as well, since they correspond to arbitrary tensor products instead of V_+, V_- .

Exempli gratia (The Hopf link - part one)

$$LT(L_D)v_{3^3} = F_5 F_3 F_2 F_1 B_3 B_2 F_5^{(2)} F_4^{(2)} F_3^{(2)} F_1^{(2)} F_2^{(3)} v_{3^2}$$



The corresponding four webs can **always** be evaluated using tableaux.

Please, fasten your seat belts!

Let's **categoryfy** everything!

A **pre-foam** is a cobordism with singular arcs between two webs. Composition consists of placing one pre-foam on **top** of the other. The following are called the **zip** and the **unzip** respectively.



They have **dots** that can move **freely** about the facet on which they belong, but we do **not** allow dot to cross singular arcs.

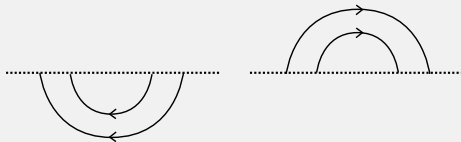
A **foam** is a formal \mathbb{C} -linear combination of isotopy classes of pre-foams modulo relations, e.g.

$$\begin{array}{c} \alpha \\ \beta \\ \delta \end{array} \begin{array}{c} \bullet \\ \bullet \\ \bullet \end{array} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \begin{array}{c} \bullet \\ \bullet \\ \bullet \end{array} = \begin{cases} 1, & (\alpha, \beta, \delta) = (1, 2, 0) \text{ or a cyclic permutation,} \\ -1, & (\alpha, \beta, \delta) = (2, 1, 0) \text{ or a cyclic permutation,} \\ 0, & \text{else.} \end{cases} \quad (\Theta)$$

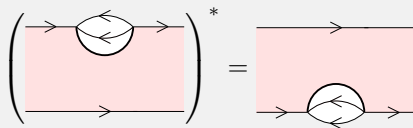
Involution on webs and foams

Definition

There is an **involution** $*$ on the webs and foams. That is



for webs and for foams



A **closed foam** is a foam from \emptyset to a closed web u^*v .

The \mathfrak{sl}_3 -foam category

Foam₃ is the **category of foams**, i.e. **objects** are webs w and **morphisms** are foams F between webs. The category is **graded** by the **q -degree**

$$\deg_q(F) = \chi(\partial F) - 2\chi(F) + 2d + b,$$

where d is the number of dots and b is the number of vertical boundary components. The **foam homology** of a closed web w is defined by

$$\mathcal{F}(w) = \mathbf{Foam}_3(\emptyset, w).$$

$\mathcal{F}(w)$ is a graded, complex vector space, whose q -dimension can be computed by the **Kuperberg bracket** (that is counting all flows on w and their weights).

A “higher” connection to knot theory

Let L_D be a link projection. Assign to it a **complex** $[[L_D]]$ in the category of formal chain complexes of **Foam**₃ locally as follows.

- $[[\text{zip}]] = 0 \rightarrow \text{)} \{2\} \xrightarrow{d} \text{X} \{3\} \rightarrow 0$.
- $[[\text{unzip}]] = 0 \rightarrow \text{X} \{-3\} \xrightarrow{d} \text{)} \{-2\} \rightarrow 0$.
- The differentials d are (un)zips.
- “Tensor” everything together.

Theorem(Khovanov)

The complex $[[L_D]]$ a link **invariant**. Moreover, it **decategorifies** to $P_3(L_D)$.

For example

$$[[\text{link}]] = 0 \longrightarrow \text{circle} \otimes \text{circle} \{2\} - \text{foam} \longrightarrow \text{zip} \{3\} \longrightarrow 0.$$

Definition(MPT)

Let $S = (s_1, \dots, s_n)$. The \mathfrak{sl}_3 -web algebra K_S is defined by

$$K_S = \bigoplus_{u,v \in B_S} {}_u K_v,$$

with

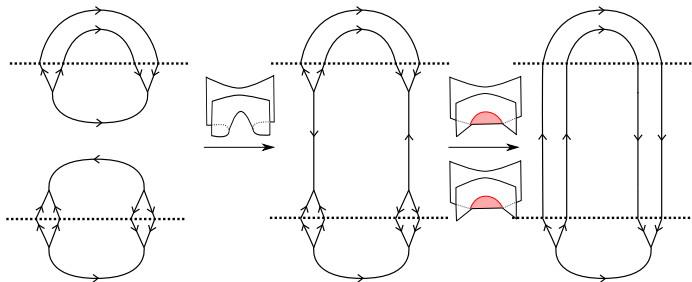
$${}_u K_v = \mathcal{F}(u^* v)\{n\}, \text{ i.e. all foams: } \emptyset \rightarrow u^* v.$$

Multiplication is defined as follows.

$${}_u K_{v_1} \otimes {}_{v_2} K_w \rightarrow {}_u K_w$$

is zero, if $v_1 \neq v_2$. If $v_1 = v_2$, use the **multiplication foam** m_v , e.g.

The \mathfrak{sl}_3 -web algebra



Theorem(s)(MPT)

The multiplication is **well-defined, associative and unital**. The multiplication foam m_v has **q -degree n** . Hence, K_S is a finite dimensional, unital and graded algebra. Moreover, it is a **graded Frobenius algebra**.

Higher representation theory

Moreover, for $n = d = 3\ell$ we define

$$W_{(3\ell)} = \bigoplus_{\mu_S \in \Lambda(n, n)_3} W_S$$

on the **level** of webs and on the **level** of foams we define

$$\mathcal{W}_{(3\ell)}^{(\rho)} = \bigoplus_{\mu_S \in \Lambda(n, n)_3} K_S - (\rho)\mathbf{Mod}_{gr}.$$

With this constructions we obtain our first **categorification** result.

Theorem(MPT)

$$K_0(\mathcal{W}_{(3\ell)}) \otimes_{\mathbb{Z}[q, q^{-1}]} \mathbb{Q}(q) \cong W_{(3\ell)} \text{ and } K_0^\oplus(\mathcal{W}_{(3\ell)}^\rho) \otimes_{\mathbb{Z}[q, q^{-1}]} \mathbb{Q}(q) \cong W_{(3\ell)}.$$

Categorification of the LT-algorithm

As a reminder, the LT-algorithm gives

$$\text{LT}(T) = x_T + \sum_{T' \prec T} \alpha_{T'}(v)x_{T'}.$$

Thus, we **need** column-strict tableaux and 3-multitableaux (for $x_{T'}$ and $\alpha_{T'}(v)$).
What do we expect to **gain**? Since Leclerc-Toffin also showed

$$b_T = \text{LT}(T) + \sum_{T'' \prec T} \beta_{T''T}(v)\text{LT}(T'')$$

we expect that we get a method to **“compute”** the projective indecomposable of K_S , since they should **decategorify** to the dual canonical basis.

A growth algorithm for foams

Definition(T)

Given a pair of a sign string and a state string (S, J) , the corresponding 3-multipartition $\vec{\lambda}$ and two Kuperberg webs $u, v \in B_S$ that extend J to f_u and f_v respectively. We define a **foam** by

$$\mathcal{F}_{\vec{T}(u_{f_u}), \vec{T}(v_{f_v})}^{\vec{\lambda}} = \underbrace{\mathcal{F}_{\sigma_u}}_{\text{Topology Idempotent}} \underbrace{e(\vec{\lambda})}_{\text{Dots}} \underbrace{d(\vec{\lambda})}_{\text{Topology}} \underbrace{\mathcal{F}_{\sigma_v}^*}_{\text{Topology}} .$$

Theorem(T)

The growth algorithm for foams is **well-defined**, the **only** input data are webs and flows on webs, works **inductively** and gives a **graded cellular basis** of K_S .

Connection to $\mathbf{U}_q(\mathfrak{sl}_d)$

Khovanov and Lauda's diagrammatic categorification of $\mathbf{U}_q(\mathfrak{sl}_d)$, denoted $\mathcal{U}(\mathfrak{sl}_d)$, is also **related** to our framework! Roughly, it consists of string diagrams of the form

$$\begin{array}{c} \nearrow \\ \nwarrow \\ \nearrow \\ \nwarrow \end{array} \begin{array}{c} \lambda \\ \lambda \end{array} : \mathcal{E}_i \mathcal{E}_j \mathbf{1}_\lambda \Rightarrow \mathcal{E}_j \mathcal{E}_i \mathbf{1}_\lambda \{(\alpha_i, \alpha_j)\}, \quad \begin{array}{c} \downarrow \\ \downarrow \\ \downarrow \\ \downarrow \end{array} \begin{array}{c} \lambda \\ \lambda \\ \lambda \\ \lambda \end{array} : \mathcal{F}_i \mathbf{1}_\lambda \Rightarrow \mathcal{F}_i \mathbf{1}_\lambda \{\alpha^{\#i}\}$$

with a weight $\lambda \in \mathbb{Z}^{n-1}$ and suitable shifts and relations like

$$\begin{array}{c} \nearrow \\ \nwarrow \\ \nearrow \\ \nwarrow \end{array} \begin{array}{c} \lambda \\ \lambda \end{array} \begin{array}{c} \bullet \\ \bullet \end{array} = \begin{array}{c} \nearrow \\ \nwarrow \\ \nearrow \\ \nwarrow \end{array} \begin{array}{c} \lambda \\ \lambda \end{array} \begin{array}{c} \bullet \\ \bullet \end{array} \quad \text{and} \quad \begin{array}{c} \nearrow \\ \nwarrow \\ \nearrow \\ \nwarrow \end{array} \begin{array}{c} \lambda \\ \lambda \end{array} \begin{array}{c} \bullet \\ \bullet \end{array} = \begin{array}{c} \nearrow \\ \nwarrow \\ \nearrow \\ \nwarrow \end{array} \begin{array}{c} \lambda \\ \lambda \end{array} \begin{array}{c} \bullet \\ \bullet \end{array}, \quad \text{if } i \neq j.$$

We define a 2-functor

$$\Psi: \mathcal{U}(\mathfrak{sl}_d) \rightarrow \mathcal{W}_{(3^\ell)}^{(\rho)}$$

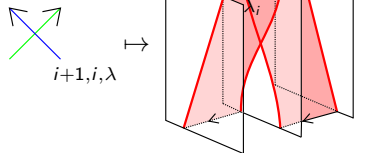
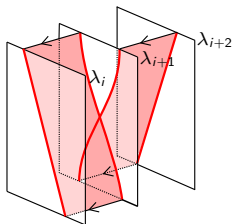
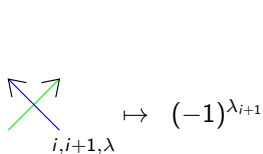
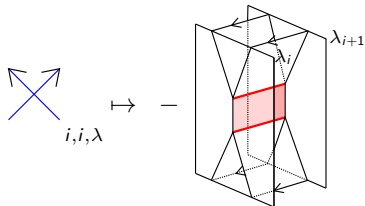
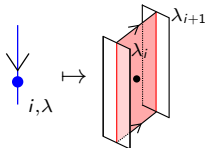
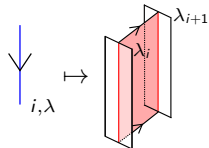
called **foamation**, in the following way.

On objects: The functor is defined by sending an \mathfrak{sl}_d -weight $\lambda = (\lambda_1, \dots, \lambda_{d-1})$ to an object $\Psi(\lambda)$ of $\mathcal{W}_{(3^\ell)}^{(\rho)}$ by

$$\Psi(\lambda) = S, \quad S = (a_1, \dots, a_\ell), \quad a_i \in \{0, 1, 2, 3\}, \quad \lambda_i = a_{i+1} - a_i, \quad \sum_{i=1}^{\ell} a_i = 3^\ell.$$

On morphisms: The functor on morphisms is by glueing the ladder webs from before on top of the \mathfrak{sl}_3 -webs in $W_{(3^\ell)}$.

On 2-cells: We define



And some others.

Theorem(MPT)

The 2-functor $\Psi: \mathcal{U}(\mathfrak{sl}_d) \rightarrow \mathcal{W}_{(3^\ell)}^{(p)}$ categorifies q -skew Howe duality.

Corollary(“Almost” directly)

We have $\psi_p([D_p^\lambda]) = b^\lambda$ under the isometry

$$\psi_p: K_0^\oplus(\mathcal{W}_{(3^\ell)}^p) \otimes_{\mathbb{Z}[q, q^{-1}]} \mathbb{Q}(q) \rightarrow W_{(3^\ell)},$$

that is projective covers D_p^λ (who give a complete list of all projective, irreducible K_S -modules) of the simple heads D^λ of the cell modules C^λ categorify the upper global crystal basis b^λ of $W_{(3^\ell)}$. In principle, the D_p^λ are computable from the extended growth algorithm.

Crossings as complexes

In order to make sense of the minus signs from before we **have to** go to the category of complexes on the categorified level. That is, define **higher** braiding operators \mathcal{B}_j (and other, similar ones)

$$\begin{array}{c} \swarrow \\ \searrow \end{array} = \left(\begin{array}{c} \begin{array}{c} 0 \\ \uparrow \\ 0 \\ \downarrow \\ 1 \end{array} \xrightarrow{\quad} \begin{array}{c} \downarrow 2 \\ \uparrow 3 \\ \downarrow 2 \end{array} \xrightarrow{\quad} \begin{array}{c} \uparrow 1 \\ 0 \\ \downarrow 0 \end{array} \\ \begin{array}{c} \downarrow 2 \\ \uparrow 3 \\ \downarrow 2 \end{array} \xrightarrow{\quad} \begin{array}{c} \uparrow 1 \\ 0 \\ \downarrow 0 \end{array} \end{array} \right) \begin{array}{c} \text{foam diagram} \\ \{ -3 \} \end{array} \left(\begin{array}{c} \begin{array}{c} 0 \\ \downarrow \\ 1 \\ \uparrow \\ 1 \end{array} \xrightarrow{\quad} \begin{array}{c} \downarrow 2 \\ \uparrow 1 \\ \downarrow 2 \end{array} \xrightarrow{\quad} \begin{array}{c} \uparrow 1 \\ 1 \\ \downarrow 0 \end{array} \\ \begin{array}{c} \downarrow 2 \\ \uparrow 1 \\ \downarrow 2 \end{array} \xrightarrow{\quad} \begin{array}{c} \uparrow 1 \\ 1 \\ \downarrow 0 \end{array} \end{array} \right) \begin{array}{c} \text{foam diagram} \\ \{ -2 \} \end{array} \end{array}$$

where the sign is **categorified** using the language of complexes and the q using degree shifts, e.g. $\{-3\}$.

Exempli gratia (The Hopf link - part two)

The Hopf link example from before will give a complex

$$\begin{array}{ccccc}
 & & F' F_4 F_3 F_3 F_2 F v_{3^2} \{5\} & & \\
 & \nearrow & \vdots & \nwarrow & \\
 & \begin{array}{c} \text{X} \\ \text{X} \end{array} : F_2 F_3 \rightarrow F_3 F_2 & & \begin{array}{c} \text{X} \\ \text{X} \end{array} : F_4 F_3 \rightarrow F_3 F_4 & \\
 F' F_4 F_3 F_2 F_3 F v_{3^2} \{4\} & & \oplus & & F' F_3 F_4 F_3 F_2 F v_{3^2} \{6\} \\
 & \searrow & \vdots & \swarrow & \\
 & \begin{array}{c} \text{X} \\ \text{X} \end{array} : F_4 F_3 \rightarrow F_3 F_4 & & \begin{array}{c} \text{X} \\ \text{X} \end{array} : F_2 F_3 \rightarrow F_3 F_2 & \\
 & & F' F_3 F_4 F_2 F_3 F v_{3^2} \{5\} & &
 \end{array}$$

that, up to some degree conventions, agrees with \mathfrak{sl}_3 -Khovanov homology of L_D . Here $F = F_5^{(2)} F_4^{(2)} F_3^{(2)} F_1^{(2)} F_2^{(3)}$ and $F' = F_5 F_4 F_3 F_1$.

Wish(T)

Since the above construction agrees (up to some normalization/shifts) with the construction of the \mathfrak{sl}_3 -link homologies, the “higher” combinatoric of tableaux should tell us how to calculate these invariants (note: this should work for **all** $n > 1$). Moreover, the **colored** \mathfrak{sl}_3 -link homologies should fit in this framework.

There is still **much** to do...

Thanks for your attention!